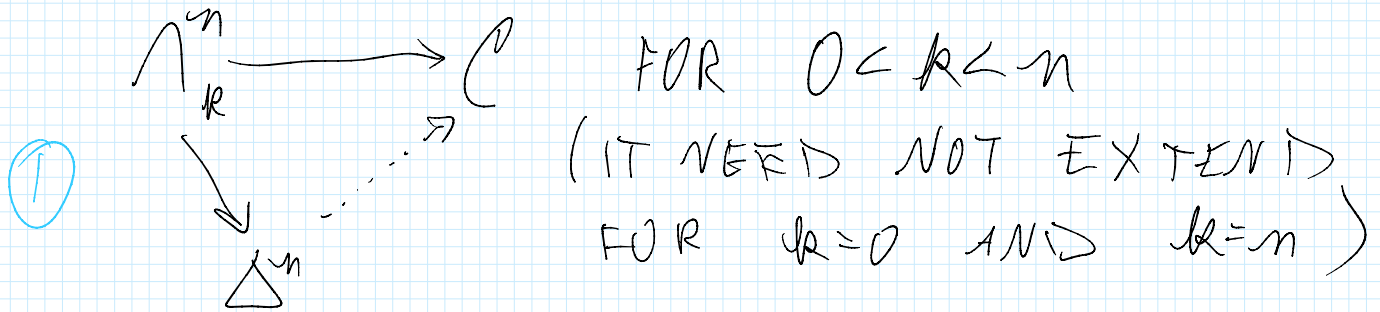


MORE ABOUT QUASI-CATEGORIES (Q-CATS)

Wednesday, April 14, 2021 10:50 AM

AKA ∞ -CATEGORIES.

RECALL A Q-CAT \mathcal{C} IS A SIMPLICIAL SET SATISFYING THE INNER HORN CONDITION



WHY IS THIS AN $(\infty, 1)$ CATEGORY?

AN ORDINARY ^{SMALL} CATEGORY \mathcal{C} HAS

A NERVE $N(\mathcal{C}) \in \mathbf{Set}_{\Delta} = \text{CAT OF SIMP SETS}$

PROP A SIMPLICIAL SET X IS THE NERVE OF A SMALL CAT \Leftrightarrow

IT IS A Q-CAT WITH UNIQUE LIFTING IN ①. THE OBJECT SET

IS X_0 AND THE MORPHISM SET IS X_1

IN A Q-CAT, COMPOSITION OF MORPHISMS IS NOT WELL DEFINED. INSTEAD

WE GET A CONTRACTIBLE SPACE OF CHOICES.

TOWARD LIMITS AND COLIMITS IN ∞ -CATS.

RECALL A TERMINAL OBJECT $*$ IN AN ORDINARY CATEGORY \mathcal{C} IS ONE THAT ADMITS A UNIQUE MORPHISM FROM ANY OBJECT X .

IN AN ∞ -CAT \mathcal{C} , WE WANT THE SIMP SET $\mathcal{C}(X, *)$ TO BE CONTRACTIBLE FOR ALL X .

RELATION BETWEEN LIMITS AND TERMINAL OBJECTS IN ORDINARY CATEGORY

EXAMPLE: EQUALIZER OF TWO MORPHISMS

$$X \begin{array}{c} \xrightarrow{a} \\ \xleftarrow{b} \end{array} Y \quad (2)$$

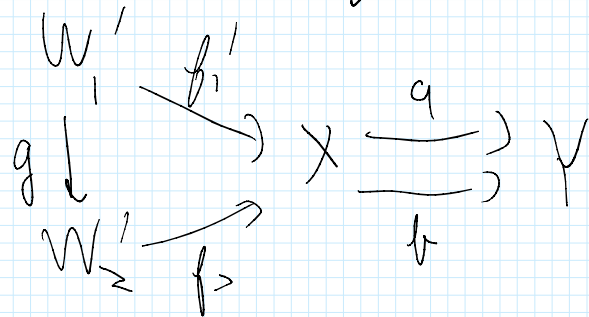
IT IS AN OBJECT W WITH A MAP TO $W \xrightarrow{f} X$ S.T. $a \circ f = b \circ f$ AND FOR ANY

$a \circ f = b \circ f$ AND FOR OTHER
 $W' \xrightarrow{b'} X \rightrightarrows Y$ WITH $a \circ f' = b \circ f'$, $\exists!$

$W' \xrightarrow{g} W$ WITH $b \circ g = b'$

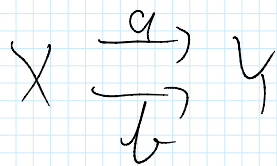
CONSIDER THE CATEGORY OF ALL
 SUCH PAIRS (W', b') AS ABOVE

WITH
 MORPHISMS



CALL THIS CATEGORY

$\mathcal{C} / X \rightrightarrows Y$ = OBJECTS IN
 \mathcal{C} LYING OVER
 THE DIAGRAM



A LIMIT W (IF IT EXISTS) IS
 A TERMINAL OBJECT IN THIS
 CATEGORY.

FOR A Q-CAT \mathcal{C} WE CAN
 DEFINE $\mathcal{C} / x \rightrightarrows y$ IN THE
 SAME WAY AND SHOW IT IS
 ALSO A Q-CAT. A LIMIT
 (IF IT EXISTS) IS A TERMINAL
 OBJECT.

DEF A Q-CAT IS COMPLETE
 (COCOMPLETE) IF ALL SMALL
 LIMITS (COLIMITS) EXIST.

HOMOTOPY IN A Q-CAT BETWEEN
 TWO MORPHISMS

$$X \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} Y \quad f \simeq g$$

ARE HOMOTOPIC IN \mathcal{C} IF

THERE IS A 2-SIMPLEX

$$\begin{array}{ccc} & X & \\ \uparrow & & \searrow g \\ 1_X & \nearrow & \\ & \mathcal{C} & \end{array}$$

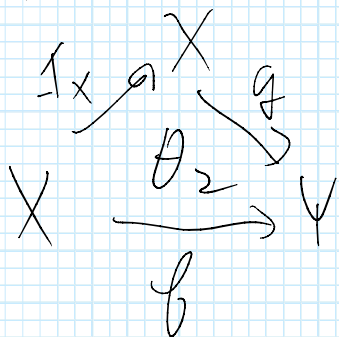
THIS IS ALSO A 2-MORPHISM

0 $\begin{array}{ccc} & & \gamma \\ & \nearrow & \searrow \\ & \theta_1 & \\ X & \xrightarrow{f} & Y \\ & & \end{array}$ THIS IS ALSO A Σ -MORPHISM $f \Rightarrow g$ IN \mathcal{C}

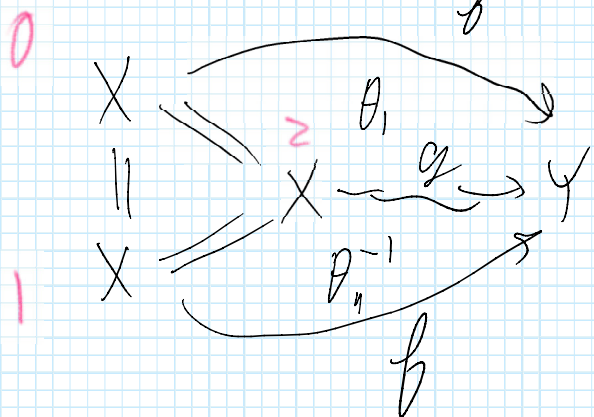
CAN SHOW THIS LEADS TO AN EQUIVALENCE RELATION ON MORPHISM $X \rightarrow Y$.

PROOF OF SYMMETRIC

NEED TO SHOW THAT IF θ_1 IS A 2-SIMPLEX, SO IS



CONSIDER A POSSIBLE 3-SIMPLEX



THE DIAGRAM DEFINES A HORN

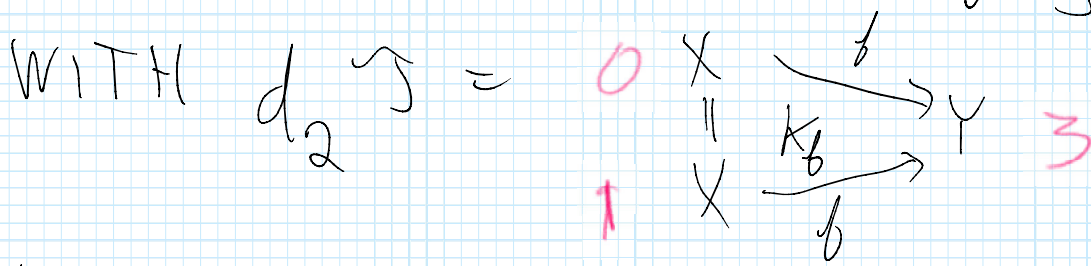
1

$$\Lambda_2^3 \longrightarrow \mathcal{C}$$

WITH FACES $(\theta_1^{-1}, \theta_1, -, X)$

$(\theta_1, \theta_1, -, X)$

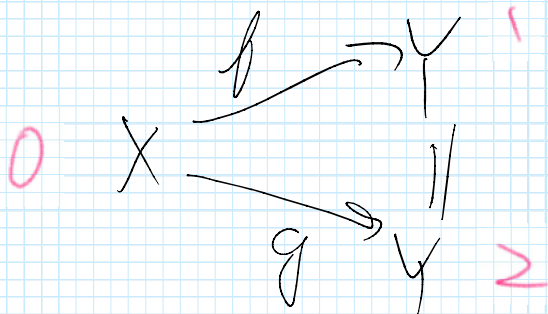
THIS MAP EXTENDS TO $\gamma: \Delta^3 \rightarrow \mathcal{C}$



THIS IMPLIES HOMOTOPY IS SYMMETRIC.

REMARK WE COULD ALSO

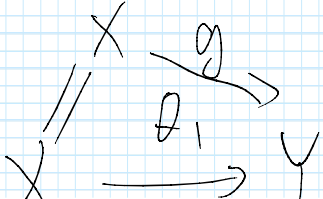
DEFINE HOMOTOPY $f \simeq g$ WITH



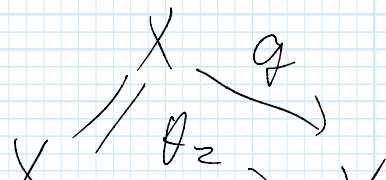
CAN SHOW THIS LEADS TO THE SAME EQUIVALENCE RELATION.

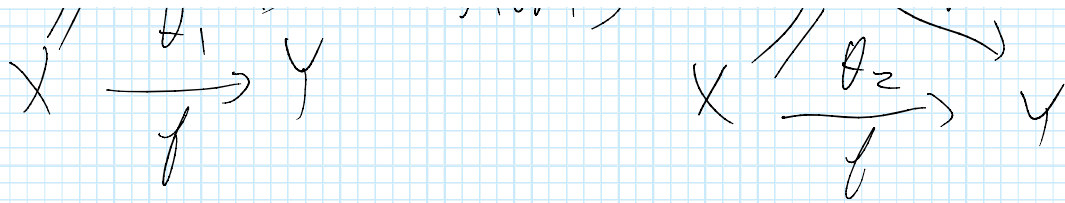
WHAT IS A 3-MORPHISM IN A Q-CAT?

SUPPOSE WE HAVE TWO 2-MORPHISM

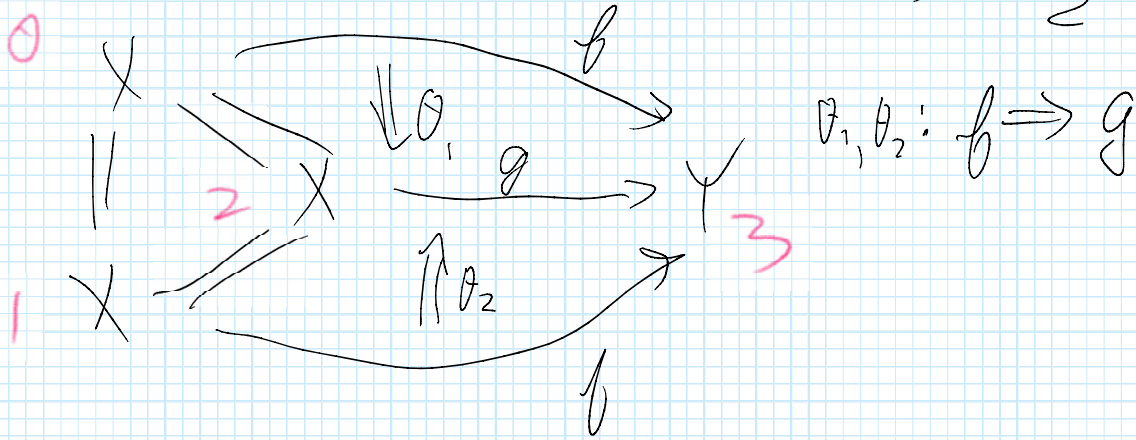


AND



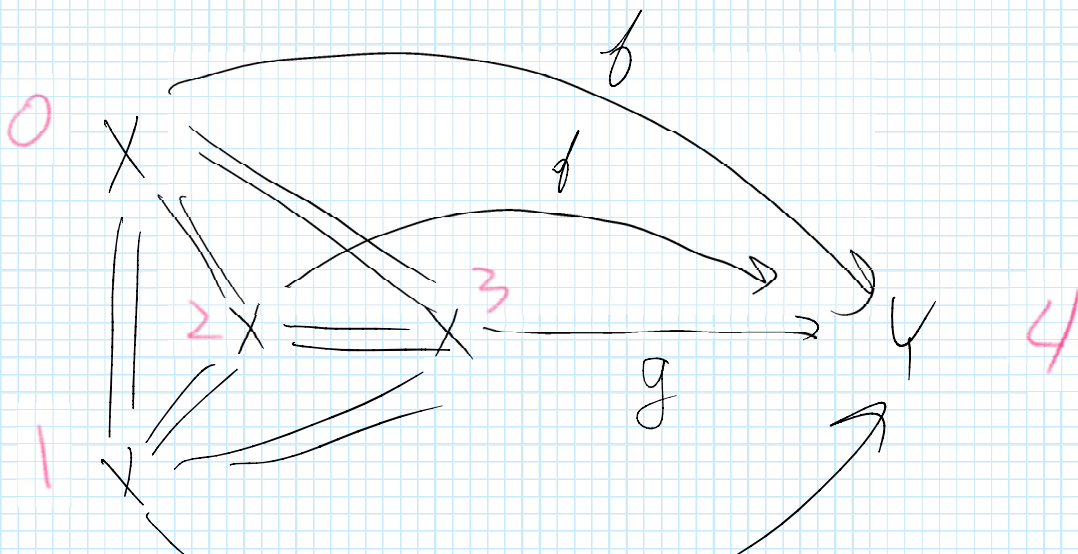


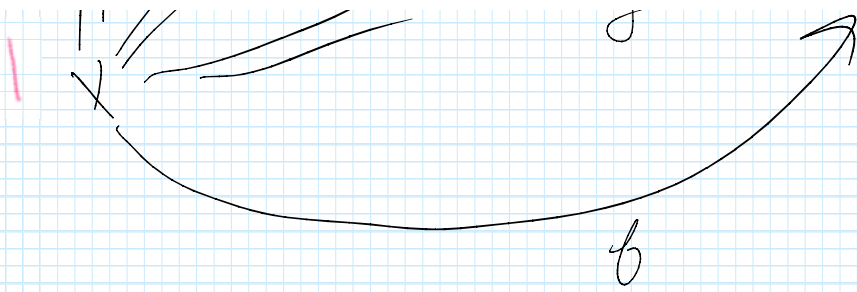
A 3-MORPHISM $S_1: \theta_1 \Rightarrow \theta_2$



WE CAN DEFINE HIGHER MORPHISMS IN A SIMILAR WAY.

FOR A 4-MORPHISM WE NEED A 4-SIMPLEX





THE VERTICES $0, 2, 3, 4$
 GIVES A 3-MORPHISM $S_1: \theta_1 \Rightarrow \theta_2$

THE VERTICES $0, 1, 3, 4$
 DEFINE ANOTHER 3-MORPHISM
 $S_2: \theta_1 \Rightarrow \theta_2$

THE WHOLE PICTURE REPRESENTS
 A 4-MORPHISM $S_1 \Rightarrow S_2$

NOTE THERE ARE 6 2-FACES,
 WITH y AS A VERTEX.

EACH REPRESENTS A 2-MORPHISM

$$0 \ 1 \ 4 \rightsquigarrow 1_f: f \Rightarrow f$$

$$0 \ 2 \ 4 \quad \text{SAME}$$

$$0 \ 3 \ 4 \quad \theta_1: f \Rightarrow g$$

$$1 \ 2 \ 4 \quad \theta_f: f \Rightarrow f$$

$$1 \ 3 \ 4 \quad \theta_2: f \Rightarrow g$$

???

1 3 4

$$\theta_2 : \theta \Rightarrow g$$

2 3 4

$$\theta_3 : \theta \Rightarrow g$$

THERE FOUR 3-FACES WITH

4 AS A VERTEX

EACH REPRESENTS

A 3-MORPHISM

1 2 3 4

0 2 3 4

0 1 3 4

0 1 2 4



WHAT ARE
THEY ?

THEY ARE EACH ONE OF THE FOLLOWING

$$S_1, S_2 : \theta_1 \Rightarrow \theta_2$$

$$I_{\theta_1} : \theta_1 \Rightarrow \theta_1$$

$$I_{\theta_2} : \theta_2 \Rightarrow \theta_2$$